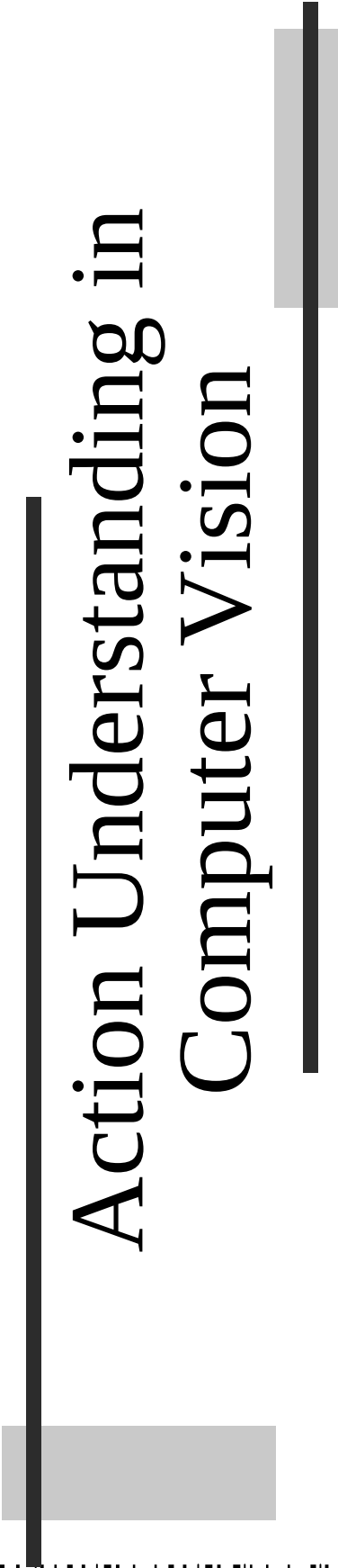


Probabilistic Context Free Grammars for Representing Action

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Action Understanding in Computer Vision

- ◆ Interpretation of basic movements
 - Sitting, walking, running ...
- ◆ Description of motion of several objects
- ◆ Recognition of gestures
- ◆ High-level event



High-level Events



- ◆ Consists of primitives
 - For car drop-off event: car-enter, car-stop, person-enter, person-exit, etc.
- ◆ Spatio-temporal structure & constraint
- ◆ Semantically defined activities
- ◆ Span extended periods of time
- ◆ Multi-object interactions



Approaches



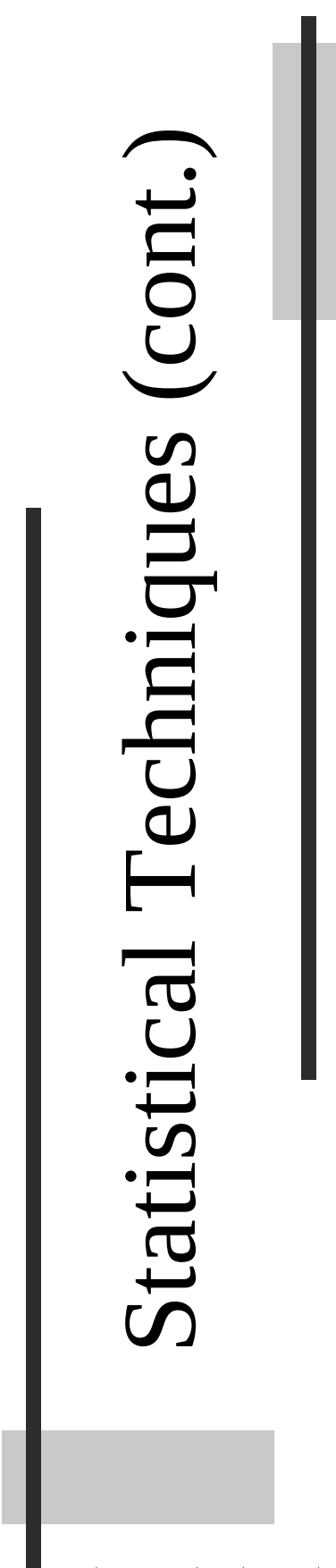
- ◆ Statistical techniques
- ◆ Syntactic techniques
- ◆ Methods that combine the two techniques



Statistical Techniques



- ◆ Classifying pattern by assuming an statistical model
 - Tennis stroke recognition
 - Gesture recognition
 - Visual language recognition
- ◆ Advantages
 - Real world data are noisy in nature (signal noise)
 - Uncertainty in observation (sensor noise)



Statistical Techniques (cont.)

- ◆ Disadvantages
 - Insufficient data
 - Semantic ambiguity
 - Temporal ambiguity
 - Known structure



Syntactic Techniques

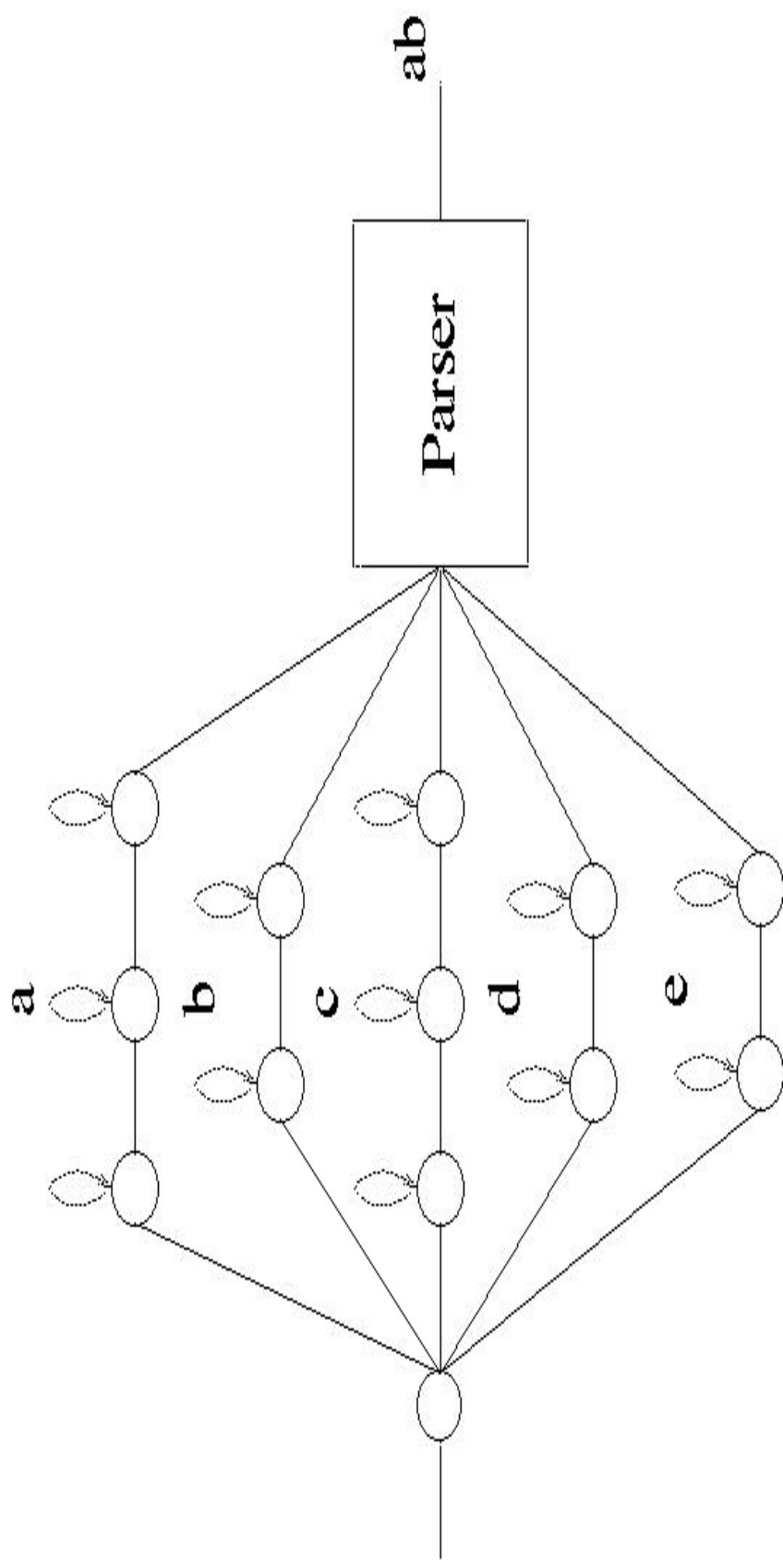


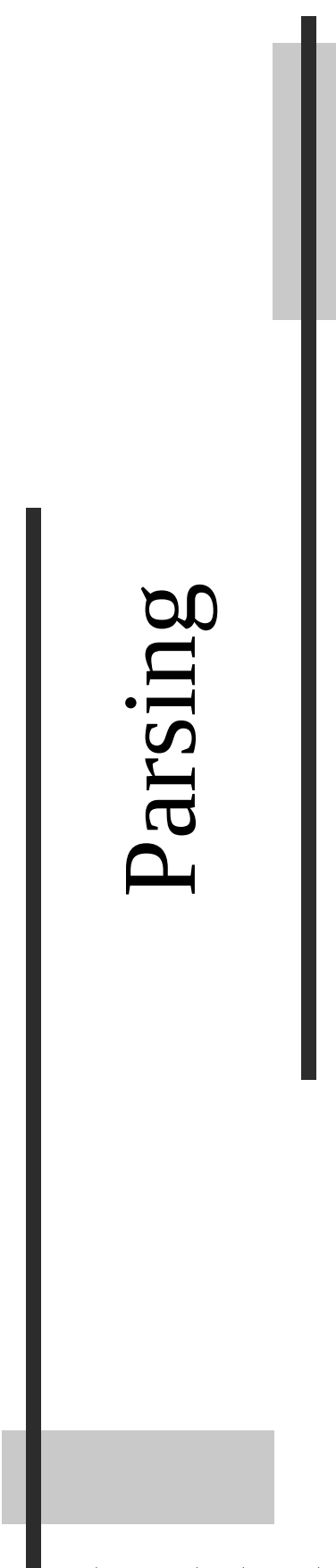
- ◆ Describe pattern structure
- ◆ Formal grammar
- ◆ Context free grammar (CFG)
- ◆ Stochastic context free grammar (SCFGs)
- ◆ Parsing

Combine the Two Techniques

- ◆ Independent primitives detection using statistical techniques
- ◆ Actions (structured primitives) recognition by syntactic techniques
 - Parsing primitives by SCFGs
 - Removing ambiguity by parsing SCFGs
 - Correcting errors (substitution, insertion, deletion) by adding SKIP rules and penalty function to SCFGs

Decoupling Primitive Detection and Primitive Structuring





Parsing

- ◆ What is Parsing?

The process of taking an input and producing some sort of structure for it. (Jurafsky & Martin)

- ◆ Structure assigned by Context Free Grammar (CFG) / Stochastic Context Free Grammar (SCFG)



Parsing Approaches



- ◆ Top-down approach
- ◆ Bottom-up approach
- ◆ Dynamic programming approach
 - Cocke-Younger-Kasami parser (CYK)
 - Graham-Harrison-Ruzzo parser (GHR)
 - Earley parser

Context Free Grammars (CFG)

- ♦ A set of non-terminal symbols N
- ♦ A set of terminal symbols Σ
- ♦ A set of productions P of form $A \rightarrow \alpha$

$$A \in N, \alpha \in (\Sigma \cup N)^*$$

- ♦ Start symbol S
- ♦ Directly derivation:
if $A \rightarrow \beta$, $\alpha, \gamma \in (\Sigma \cup N)^*$,
 $\alpha A \gamma \Rightarrow \alpha \beta \gamma$

Context Free Grammar (Cont.)

- ◆ Derivation: $\alpha_1 \rightarrow \alpha_2, \alpha_2 \rightarrow \alpha_3, \dots, \alpha_{m-1} \rightarrow \alpha_m,$

*

$$\alpha_1 \Rightarrow \alpha_m$$

- ◆ A language generated by a grammar G

$$\{L_G = w \mid w \in \Sigma^*, S \xRightarrow{*} w\}$$

Stochastic Context Free Grammar (SCFG)

- ◆ Modify production as: $A \rightarrow \lambda[p]$
- ◆ Where $p = P(A \rightarrow \lambda \mid A)$ is the rule probability of the production $A \rightarrow \lambda$ from a Context-Free Grammars (CFGs)
- ◆ Rules are conditionally independent

$$\begin{aligned} P(A \rightarrow \lambda, B \rightarrow \beta, C \rightarrow \gamma \mid A, B, C) \\ = P(A \rightarrow \lambda \mid A) \cdot P(B \rightarrow \beta \mid B) \cdot P(C \rightarrow \gamma \mid C) \end{aligned}$$

Earley Parsing Algorithm

- ◆ A set of states for each position in the input
- ◆ Dot denotes the current input position
- ◆ A state with the dot at the right most position is a complete state
- ◆ A state produced by prediction is a predicted state
- ◆ A state produced by completion is a completed state

Earley Parsing Algorithm

(cont.)

♦ A State: $i : X_k \rightarrow \lambda . Y\mu$

$\begin{array}{c} \text{a b c d e f g h s d} \dots \\ \uparrow \quad \downarrow \\ \quad \quad \quad \text{i} \quad \quad \quad \text{k} \end{array}$

♦ Prediction:

$$\left\{ \begin{array}{l} i : X_k \rightarrow \lambda . Y\mu, \\ Y \rightarrow \nu, \end{array} \right. \Rightarrow i : Y_i \rightarrow . \nu$$

Earley Parser (cont.)

- ◆ Scanning: $i : X_k \rightarrow \lambda . a \mu \Rightarrow i+1 : X_k \rightarrow \lambda a . \mu$
- ◆ Completion:

$$\begin{cases} j : X_k \rightarrow \lambda . Y \mu, \\ i : Y_j \rightarrow \nu ., \end{cases} \Rightarrow i : X_k \rightarrow \lambda Y . \mu$$

a b c d e f g h s d r w v . . .

↑_k ↓_j ↑_i

An Example

$S \rightarrow NP VP$
 $S \rightarrow Aux NP VP$
 $S \rightarrow VP$
 $NP \rightarrow Det Noun$
 $VP \rightarrow Verb$

state set (0)

$0:0_0 \rightarrow \cdot S$
predicted
 $0:S_0 \rightarrow \cdot NP VP$
 $0:NP_0 \rightarrow \cdot Det Noun$
 $0:NP_0 \rightarrow \cdot ProperNoun$
 $0:S_0 \rightarrow \cdot Aux NP VP$
 $0:S_0 \rightarrow \cdot VP$
 $0:VP_0 \rightarrow \cdot Verb$
 $0:VP_0 \rightarrow \cdot Verb NP$

$Verb \rightarrow book$
 $Det \rightarrow that$
 $Noun \rightarrow flight$
 $NP \rightarrow ProperNoun$
 $VP \rightarrow Verb VP$

(1) Book

scanned
 $1:Verb_0 \rightarrow book \cdot$
completed
 $1:VP_0 \rightarrow Verb \cdot$
 $1:S_0 \rightarrow VP \cdot$
 $1:VP_0 \rightarrow Verb \cdot NP$
predicted
 $1:NP_1 \rightarrow \cdot Det Noun$
 $1:NP_1 \rightarrow \cdot ProperNoun$

(2) that

scanned
 $2:Det_1 \rightarrow that \cdot$
completed
 $2:NP_1 \rightarrow Det \cdot Noun$

(3) flight

scanned
 $3:Noun_2 \rightarrow flight \cdot$
completed
 $3:NP_1 \rightarrow Det Noun \cdot$
 $3:VP_0 \rightarrow Verb NP \cdot$
 $3:S_0 \rightarrow VP \cdot$

Earley-Stolcke Parser (1)

- ◆ A state $i : X_k \rightarrow \lambda \cdot Y\mu [\alpha, \gamma]$
- ◆ Forward probability
 $\alpha = \sum_{path} P(\text{path of length } i \text{ that ends in state } i : X_k \rightarrow \lambda \cdot Y\mu)$
- ◆ Inner probability
 $\gamma = \sum_{path} P(\text{path of length } i - k \text{ that start in state } k : X_k \rightarrow \cdot \lambda Y\mu \text{ and end in state } i : X_k \rightarrow \lambda \cdot Y\mu)$
- ◆ Earley path: a sequence of states needed to reach the current state
- ◆ Length of path: number of scanning states

Earley-Stolcke Parser (2)

- ◆ Prediction

$$\left\{ \begin{array}{l} i : X_k \rightarrow \lambda \cdot Z\mu[\alpha, \gamma] \\ Y \rightarrow \nu, \end{array} \right. \Rightarrow i : Y_i \rightarrow \bullet \nu[\alpha', \gamma']$$

where

$$\alpha' = \sum_{\lambda, \mu} \alpha(i : X_k \rightarrow \lambda \cdot Z\mu) \cdot R_L(Z, Y) \cdot P(Y \rightarrow \nu)$$

$$\gamma' = P(Y \rightarrow \nu)$$

Compute R_L

- ◆ Left-recursion in grammar
- ◆ Possibly infinite prediction loop that accumulate probability computation
- ◆ Example: $A \rightarrow Aa, A \rightarrow a$
- ◆ Left Corner relation: $X \rightarrow_L Y$, iff $\forall X \rightarrow Y\lambda$
- ◆
$$R_L(Z, Y) = P_L(Z \Rightarrow^0 Y) + P_L(Z \Rightarrow^1 Y) + \dots$$
- ◆
$$P_L(Z \Rightarrow^k Y) = P_L(Z \Rightarrow Y_1 \Rightarrow Y_2 \Rightarrow \dots \Rightarrow Y_{k-1} \Rightarrow Y)$$

Compute R_L (cont.)

- ◆ Matrix form

$$R_L = P_L^0 + P_L^1 + \dots = \sum_{k=0}^{\infty} (I - P_L)^{-1}$$

- ◆ Computed once for the grammar, and used at each iteration of the prediction step

Earley-Stolcke Parser (3)

- ◆ Scanning

$$i:X_k \rightarrow \lambda . a\mu [\alpha, \gamma] \Rightarrow i+1:X_k \rightarrow \lambda a . \mu [\alpha, \gamma]$$

- ◆ Completion

$$\begin{cases} j:X_k \rightarrow \lambda . Z\mu[\alpha, \gamma], \\ i:Y_j \rightarrow \nu . [\alpha', \gamma'], \end{cases} \Rightarrow i:X_k \rightarrow \lambda Z . \mu[\alpha', \gamma']$$

where

$$\alpha' = \sum_{\nu} \alpha(j:X_k \rightarrow \lambda . Z\mu) \cdot R_U(Z, Y) \cdot \gamma''(i:Y_j \rightarrow \nu.)$$

$$\gamma' = \sum_{\nu} \gamma(j:X_k \rightarrow \lambda . Z\mu) \cdot R_U(Z, Y) \cdot \gamma''(i:Y_j \rightarrow \nu.)$$

Compute R_U

- ◆ Unit production: $X \rightarrow_U Y$, iff $X \rightarrow Y$
- ◆ Infinite completion by unit production
 - e. g. $A \rightarrow B, A \rightarrow a, B \rightarrow A$
- ◆ Unit production relation matrix P_U
- ◆ Similarly as computing R_L in prediction step

$$R_U = P_U^0 + P_U^1 + \dots = \sum_{k=0}^{\infty} (I - P_U)^{-1}$$

Uncertainty in the Input

- ◆ Source of the input symbols is probabilistic
- ◆ Modify scanning of the Earley-Stolcke parser

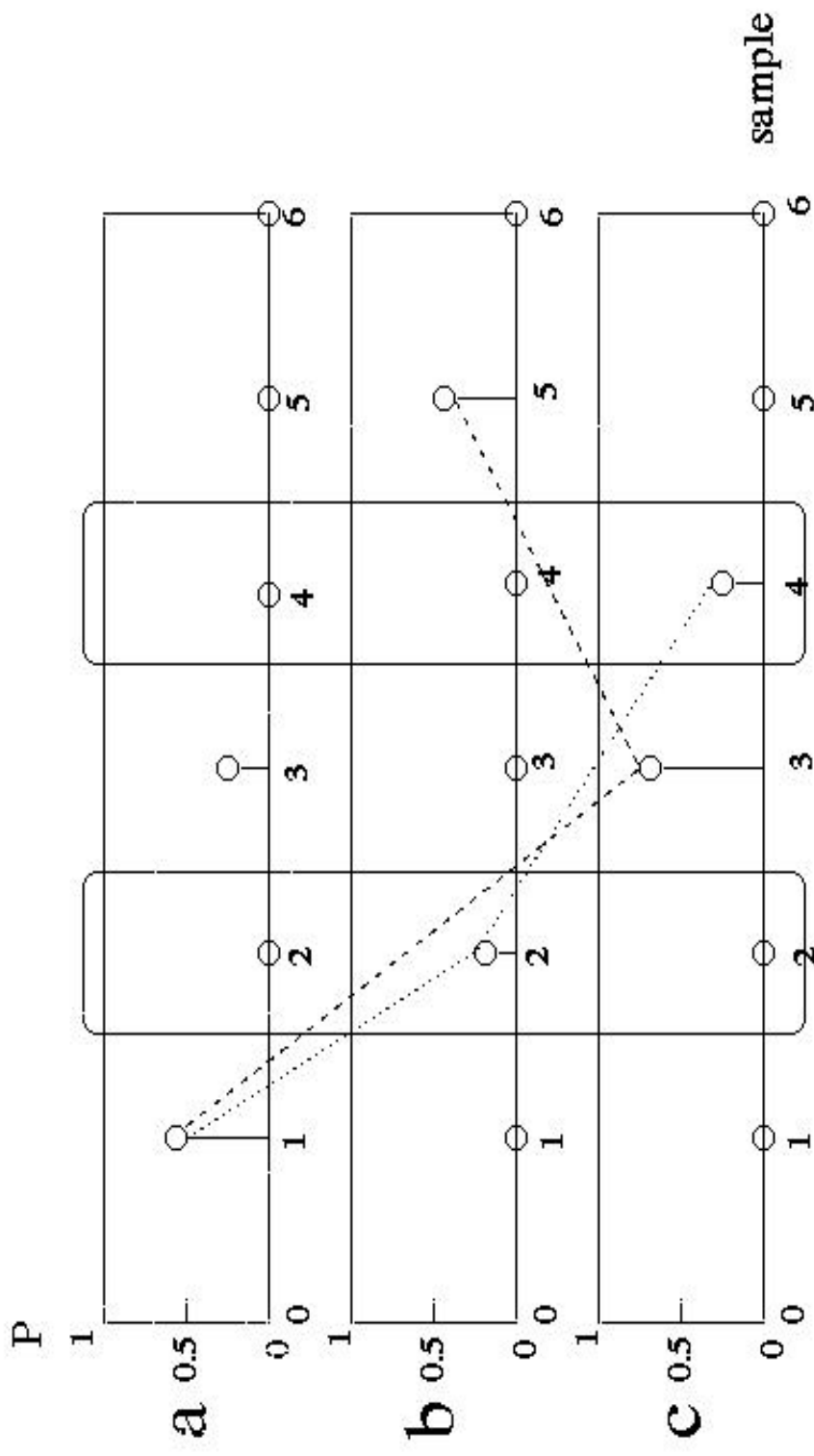
$$\left\{ \begin{array}{l} i : X_k \rightarrow \lambda . a \mu \ [\alpha, \gamma] \Rightarrow i+1 : X_k \rightarrow \lambda a . \mu \ [\alpha', \gamma'] \\ \forall a, \text{ s.t. } P(a) > 0 \end{array} \right.$$

$$\alpha' = \alpha (i : X_k \rightarrow \lambda . \alpha \mu) \cdot P(a)$$

$$\gamma' = \gamma (i : X_k \rightarrow \lambda . \alpha \mu) \cdot P(a)$$

- ◆ Address the substitution error

An Example for a Grammar

$$A \rightarrow abc \mid acb$$


Insertion and Deletion

- ◆ Use a robust form grammar $\hat{G}$ of G
 $G: \Rightarrow \hat{G}$

$$A \rightarrow bC \quad A \rightarrow \hat{B}C$$

- ◆ $\hat{B} \rightarrow b \mid \text{SKIP } b \mid b \text{ SKIP}$
- ◆ $\text{SKIP} \rightarrow b \mid c \mid \dots \mid b \text{ SKIP} \mid \dots$
- ◆ Includes all repetitions of all terminals
- ◆ Set $P(\text{SKIP rule})$ small
- ◆ Penalize derivation consuming less terminals

Enforcing Consistency (1)

- ◆ Types of consistency
 - Temporal consistency
 - Spatial consistency
 - Object identity consistency
- ◆ Add 2 vector valued state variables
 - l : low mark
 - h : high mark
- ◆ Containing the data for computing distance penalty between two joining states

Enforcing Consistency (2)

- ◆ Prediction

$$\left\{ \begin{array}{l} i : X_k \rightarrow \lambda . Z\mu[l, h] \\ Y \rightarrow \nu, \end{array} \right. \Rightarrow i : Y_i \rightarrow \bullet . \nu[S_t, S_t]$$

- ◆ Scanning

$$\begin{array}{l} i : X_k \rightarrow \lambda . a\mu[l, h] \Rightarrow \\ \left\{ \begin{array}{l} i+1 : X_k \rightarrow \lambda a . \mu[l_a, h_a] \text{ if } \lambda = \varepsilon \\ i+1 : X_k \rightarrow \lambda a . \mu[l, h_a] \text{ else} \end{array} \right. \end{array}$$

Enforcing Consistency (3)

- ◆ Completion

$$\alpha' = f(d) \sum_{\nu} \alpha(j : X_k \rightarrow \lambda \cdot Z\mu) \cdot R_U(Z, Y) \cdot \gamma''(i : Y_j \rightarrow \nu \cdot)$$

$$\gamma' = f(d) \sum_{\nu} \gamma(j : X_k \rightarrow \lambda \cdot Z\mu) \cdot R_U(Z, Y) \cdot \gamma''(i : Y_j \rightarrow \nu \cdot)$$

- ◆ $f(d)$: distance penalty function
- ◆ Computed based on high mark of completed state and low mark of completing state

Choice of $f(d)$

- ◆ Sever penalty: step function

- e. g.

$$f(d) = \begin{cases} 0, & d < 0 \\ 1, & \text{else} \end{cases}$$

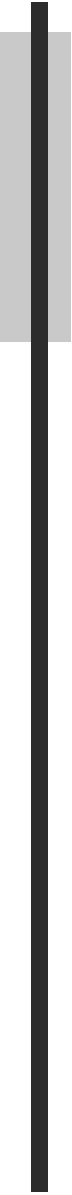
- ◆ Softer penalty: exponential function

- e. g.

$$f(d) = Ce^{-\frac{d^2}{\theta}}$$



Application: Video Surveillance of Parking Lot



- ◆ Outdoor environment – occlusions and lighting change
- ◆ Static cameras
- ◆ Real-time performance
- ◆ Labeling activities and person-vehicle interactions in a parking lot
- ◆ Handling simultaneous events



Known Structure, Uncertain Elements

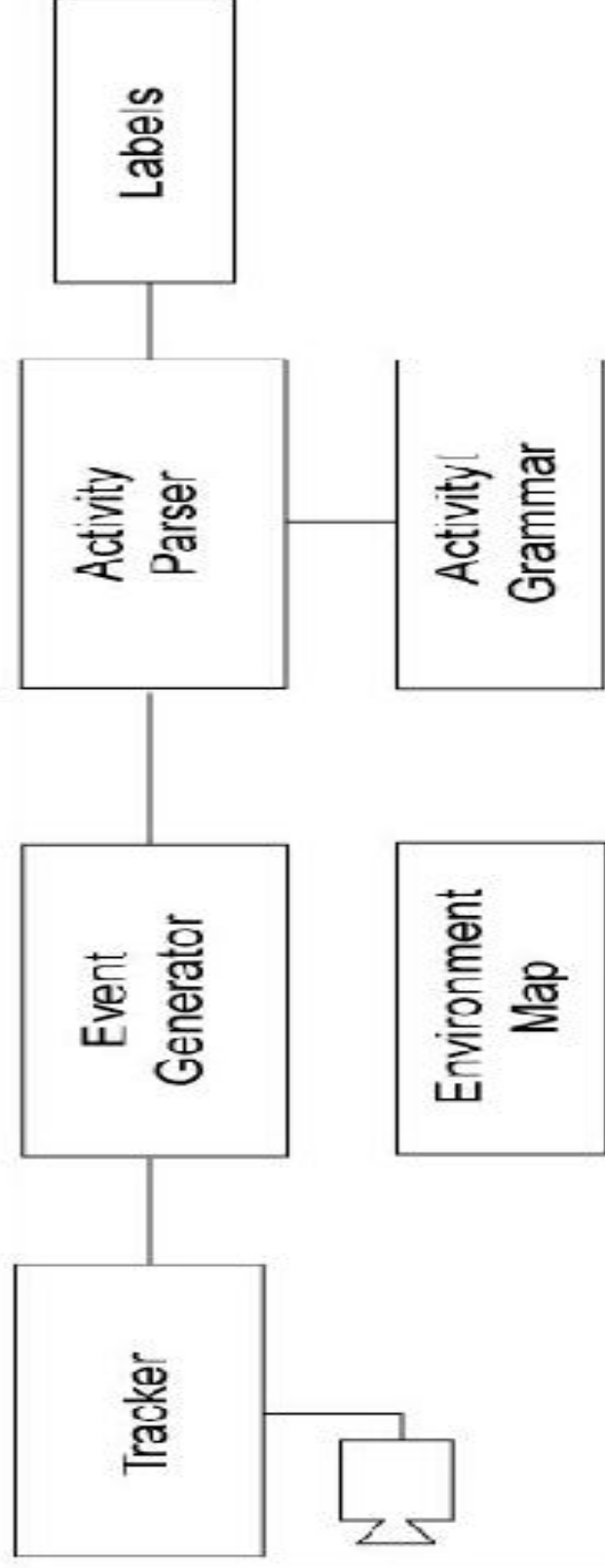


- ◆ Activities as sequences of primitives represented by SCFG
 - Car drop-off, car pick-up
 - Dancing
- ◆ Input primitives are uncertain
 - Uncertain observation of primitive
 - Noisy symbols

Approach

- ◆ First detect primitives using statistical method
 - Tracker
 - Event generator
- ◆ Then Recognize activity by parsing input stream of uncertain primitives (partial tracks) by an SCFG parser

System Overview



System Overview (cont.)

- ◆ Tracker
 - Assign identity to the moving objects
 - Collects the trajectory data into partial tracks
- ◆ Event generator
 - Maps partial tracks onto predetermined set of events
- ◆ Parser
 - Labels sequences of events by parsing using a SCFG
 - Enforce consistency constraint



Tracker



- ◆ Object found
- ◆ Assign a unique ID
- ◆ Track changes in objects' appearance, position, velocity
- ◆ Based on the data, assign each object a class label (e.g. a car or a person)
- ◆ Object lost
- ◆ Object exit

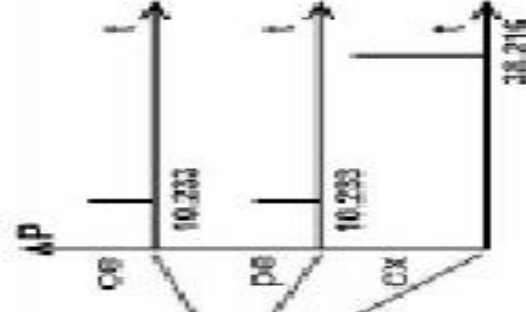
Event Generator

- ◆ Based on data from tracker
 - Object-enter
 - Object-found
 - Object-exit
 - Object-lost
 - Object-stopped
- ◆ Initially, tracker can not figure out class label,
- ◆ When object exit, tracker has enough information to assign a class label to the object

An Example of Generating Events



Event	Likelihood	x	y	dx	dy	time
car-enter	0.5	0.454	1	-0.01	0.05	10.233
person-enter	0.5	0.454	1	-0.01	0.05	10.233
car-exit	1	1	0.784	0.1	0.1	38.216



Parsing Events

$$\begin{array}{l} X \rightarrow Zc \\ Z \rightarrow ab \end{array}$$

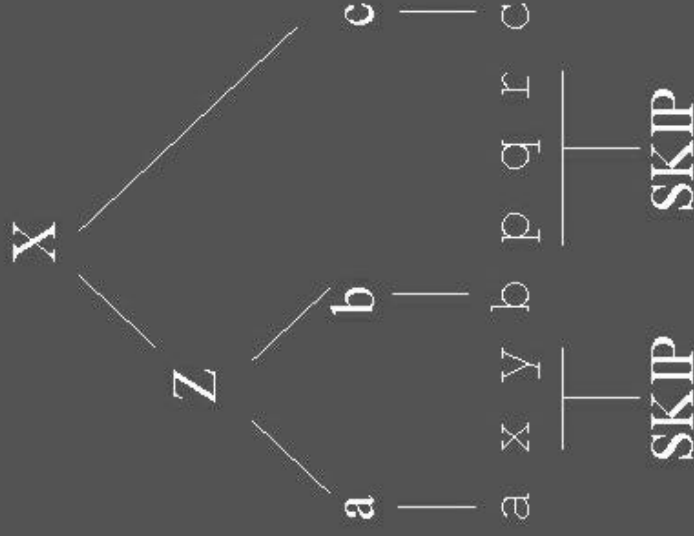
- production rules (states)

X - target non-terminal (label)

Z - intermediate non-terminal

- input stream (tracking events)

- noise rules



For the production X, events a, b and c should be consistent

Sample Stochastic Context-Free Grammar

Partial Listing:

```
TRACK      -> CAR-TRACK [0.50]
           | PERSON-TRACK [0.50]

CAR-TRACK  -> CAR-THROUGH [0.25]
           | CAR-PICKUP [0.25]
           | CAR-OUT [0.25]
           | CAR-DROP [0.25]

CAR-THROUGH . . .

. . .

CAR-EXIT   -> car-exit [0.70]
           | SKIP car-exit [0.30]
```

Non-terminals -
semantically significant
groups of events

Terminals - individual
events

Rule probabilities

SKIP-rule - noise symbol

Tracker and Event Generator

Data for Parser States

- ◆ Tracker event generator provides data for “low mark” and “high mark” of parser states

$$l = \begin{pmatrix} f_l \\ t_l \\ x_l \\ y_l \\ dx_l \\ dy_l \end{pmatrix} \quad h = \begin{pmatrix} f_h \\ t_h \\ x_h \\ y_h \\ dx_h \\ dy_h \end{pmatrix}$$

f: frame number
t: timing stamp
(x,y): location
(dx,dy): velocity

Distance Penalty Function

- ♦ t_1, r_1, dr_1 : high mark data of state being completed
- ♦ t_2, r_2 : low mark of the completing state

$$f(t_1, t_2, r_p, r_2) = \begin{cases} 0, & \text{if } (t_2 - t_1) < 0 \\ \exp\left(\frac{(r_2 - r_p)^T (r_2 - r_p)}{\theta}\right), & \text{else} \end{cases}$$

- ♦ Where $r_p = r_1 + dr_1(t_2 - t_1)$ is the predicted position of the object at time t_2

An Example

Completion - check if Z can be used as a supporting evidence for X :

$$i: X_k \rightarrow \lambda. Z \mu \quad [\alpha_1, \gamma_1, l_1, h_1]$$

$?\uparrow$

$$j: Z_i \rightarrow \eta. \quad [\alpha_2, \gamma_2, l_2, h_2]$$

Legend:

$$\mathbf{r} = (x, y)$$

$$d\mathbf{r} = (dx, dy)$$

X Z t
reject

X Z t
accept

X Z t
penalize



Predict new position:

$$\mathbf{r}_p = \mathbf{r}_1 + d\mathbf{r}_1(t_2 - t_1)$$

Penalize:

$$f(\mathbf{r}_p, \mathbf{r}_2) = \begin{cases} 0, & \text{if } (t_2 - t_1) < 0 \\ \exp\left(\frac{(\mathbf{r}_2 - \mathbf{r}_p)^T (\mathbf{r}_2 - \mathbf{r}_p)}{\theta}\right) & \end{cases}$$

Events Data for Drive-In and Drop-Off Activities

DROP-OFF							DRIVE-IN						
Event	UID	Avg. Size	Class	P	x	y	t	frame					
ENTER	724	0.122553	0	0.5	0.450094	0.938069	917907137.8	1906					
ENTER	665	0.046437	1	0.5	0.6107	0.94674	917907122.5	1799					
PERSON-LEAVE	665	0.045869	1	0.997846	0.648089	0.98855	917907142.7	1938					
STOPPED	724		0	0.995784	0.348569	0.345513	917907146.5	1964					
ENTER	780	0.034293	1	0.5	0.74188	0.980292	917907151.3	1998					
ENTER	790	0.069093	0	0.5	0.814565	0.032611	917907153.4	2012					
FOUND	787	0.033573	1	0.5	0.297585	0.357887	917907153.1	2010					
CAR-LEAVE	790	0.061263	0	0.997285	0.975971	0.211984	917907155.3	2025					
PERSON-LEAVE	780	0.038616	1	0.999923	0.974494	0.865237	917907158.6	2047					
PERSON-LEAVE	787	0.032045	1	0.999997	0.296519	0.183704	917907158.7	2048					
ENTER	813	0.034776	1	0.5	0.012821	0.348379	917907160.9	2063					
ENTER	816	0.093513	0	0.5	0.960425	0.793899	917907161.9	2070					
CAR-LEAVE	724	0.097374	0	0.993211	0.972272	0.693728	917907165.2	2091					
CAR-LEAVE	816	0.089424	0	0.99023	0.693699	0.990798	917907165.2	2091					

Vedio Frame Illustration



Person passed through



Person drove in



Person drop off



Car passed through

